

# Digraphic Topology On Directed Edges

Hussein A. Neamah<sup>1</sup> and Khalid Sh. Al’Dzhabri<sup>2</sup>

<sup>1</sup>Department of Mathematics, University of Al-Qadisiyah , College of Education, Iraq , Al Diwaniyah  
[edu-math.post6@qu.edu.iq](mailto:edu-math.post6@qu.edu.iq)

<sup>2</sup>Department of Mathematics, University of Al-Qadisiyah , College of Education, Iraq , Al Diwaniyah  
[khalid.aljabrimath@qu.edu.iq](mailto:khalid.aljabrimath@qu.edu.iq)

**Abstract:** In this paper, we study the digraphic topology  $\tau_{ID}$  for a directed edges of a digraph. We give some properties of this topology, in particular we prove that  $\tau_{ID}$  is an Alexandroff topology and when two digraphs are isomorphic, their digraphic topologies will be homeomorphic. We give some properties matching digraphs and homeomorphic topology spaces. Finally, we investigate the connectedness of this topology and some relations between the connectedness of the digraph and the topology  $\tau_{ID}$ .

**Keywords—** digraph, topology, Alexandroff topology and separation axioms.

## 1.Introduction:

Topological structures are mathematical models that can be used to analyze data without the concept of distance. Topological structures, in our opinion, are a crucial adjustment for the extraction and processing of knowledge [2]. This publication provides a few topological fundamentals that are pertinent to our study. One of the most crucial structures in discrete mathematics is the graph [1]. Two observations explain their pervasiveness. Graphs are mathematically elegant, to start, from a theoretical standpoint. Although a graph merely has a set of vertices and a relationship between pairs of vertices, it is a simple structure, yet graph theory is a vast and diverse field of study. This is partially because graphs can be thought of as topological spaces, combinatorial objects, and many other mathematical structures in addition to being relational structures [1]. This brings us to our second argument about the significance of graphs: many ideas may be abstractly represented by graphs[3], which makes them very helpful in real-world applications. Several earlier studies on the subject of topological graphs we can see in [4-11]. In this paper we discuss a new method to generate topology  $\tau_{ID}$  on graph by using new method of taking neighborhood is determining a vertex on the digraph and calculate each vertex and its edges indgree of it and we defined  $S_{ID} = \{ \overrightarrow{E_v} | v \in V \}$ , where  $E_v$  is the set of all edges indgree to  $v$ , we have  $E = \bigcup_{v \in V} \overrightarrow{E_v}$ . hence  $S_{ID}$  forms a subbasis for a topology  $\tau_{ID}$  on  $E$ , called digraphic topology,(briefly digtopology) of  $D$ .

## 2.PRELIMINARIES:

In this work , some basic notions of graph theory [1-2], and topology [2] are presented. A graph (resp., directed graph or digraph)  $D = (V, E)$  consists of a vertex set  $V$  and an edge set  $E$  of unordered (resp., ordered) pairs of elements of  $V$ . To avoid ambiguities, we assume that the vertex and edge sets are disjoint. We say that two vertices  $v$  and  $w$  of a graph (resp., digraph  $D$ ) are adjacent if there is an edge of the form  $vw$  (resp.,  $\overrightarrow{vw}$  or  $\overleftarrow{vw}$ ) joining them, and the vertices  $v$  and  $w$  are then incident with such an edge. A subdigraph  $H$  of a digraph  $D$  is a digraph, each of whose vertices belong to  $V$  and each

of whose edges belong to  $E$ . The degree of a vertex  $v$  of  $D$  is the number of edges incident with  $v$ , and written  $\text{deg}(v)$ . A vertex of degree zero is an isolated vertex. In digraph, the outdegree, of a vertex  $v$  of  $D$  is the number of edges of the form  $\overrightarrow{vw}$  and denoted by  $d^+(v)$ , similarly, the indegree of a vertex  $v$  of  $D$  is the number of edges of the form  $\overleftarrow{vw}$ , and denoted by  $d^-(v)$ . A vertex of out-degree and in-degree are zero is an isolated vertex . A topology  $\tau$  on a set  $X$  is a combination of subset of  $X$ , called open, such that the union of the member of any subset of  $\tau$  is a member of  $\tau$ , the intersection of the members of any finite subset of  $\tau$  is a member of  $\tau$ , and both empty set and  $X$  are in  $\tau$  and the ordered pair  $(X, \tau)$  is called topological space. The topology  $\tau = P(X)$  on  $X$  is called discrete topology while the topology  $\tau = \{X, \emptyset\}$  on  $X$  is called indiscrete topology. A topology in which arbitrary intersection of open set is open called Alexandroff space.

## 3.DIGRAPHIC TOPOLOGY.

In this section, we introduce our new subbasis family to generated a topology on the set of edges  $E$  of a digraph  $D = (V, E)$ .

**Definition 3.1:** Let  $D = (V, E)$  be a digraph, we defined  $S_{ID} = \{ \overrightarrow{E_v} | v \in V \}$ , where  $E_v$  is the set of all edges indgree to  $v$ , we have  $E = \bigcup_{v \in V} \overrightarrow{E_v}$ . hence  $S_{ID}$  forms a subbasis for a topology  $\tau_{ID}$  on  $E$ , called digraphic topology,(briefly digtopology) of  $D$ .

**Theorem 3.2:** Let  $D = (V, E)$  be digraph then  $(E, \tau_{ID})$  is topological space .

**prove :** we will prove that  $\tau_{ID}$  is topological graph ,

- 1) Since  $E = \bigcup_{i \in I} W_i$  where  $W_i \in \beta_{ID}$  such that  $\beta_{ID}$  is a basis for a topological graph  $\tau_{ID}$ , then  $W_i = \bigcap_{i=1}^n S_i$  where  $S_i \in S_{ID}$  and  $S_i = \overrightarrow{E_{v_i}}$ ,  $v_i \in V$ . Then  $E = \bigcup_{i \in I} (\overrightarrow{E_{v_i}})$  and so  $E \in \tau_{ID}$ . Also  $\emptyset \in \tau_{ID}$  by complement of  $E$ .
- 2) Let  $U_i \in \tau_{ID}$ ,  $U_i = \bigcup_{i \in I} W_i$  where  $W_i \in \beta_{ID}$ ,  $W_i = \bigcap_{i=1}^n S_i$  where  $S_i \in S_{ID}$ ,  $S_i = \overrightarrow{E_{v_i}}$ ,  $v_i \in V$ , then  $U_i =$

$U_{i \in I}(\bigcap_{j=1}^n \overrightarrow{E_{v_i}})$  so  $U_i = U_{i \in I}(\bigcap_{j=1}^n \overrightarrow{E_{v_i}}) \in \tau_{ID}$ .

Therefore  $U_{i \in I} U_i \in \tau_{ID}$ .

3) Let  $A_i, B_i \in \tau_{ID}$ ,  $A_i = U_{i \in I} W_i$  where  $W_i \in \beta_{ID}$ ,  $W_i = \bigcap_{i=1}^n S_i$  where  $S_i \in S_{ID}$ ,  $S_i = \overrightarrow{E_{v_i}}$ ,  $v_i \in V$  then  $A_i = U_{i \in I}(\bigcap_{j=1}^n \overrightarrow{E_{v_i}})$ ,  $B_i = U_{i \in I} U_i$  where  $U_i \in \beta_{ID}$ ,  $U_i = \bigcap_{i=1}^n S_i$  where  $S_i \in S_{ID}$ ,  $S_i = \overrightarrow{E_{v_i}}$ ,  $v_i \in V$  then  $B_i = U_{i \in I}(\bigcap_{j=1}^n \overrightarrow{E_{v_i}})$  then:

- i) If there are no element in intersection i.e.,  $A_i \cap B_i = \emptyset$  since  $\emptyset \in \tau_{ID}$  then  $A_i \cap B_i \in \tau_{ID}$ .
- ii) If there exist element in intersection  $U_i \cap W_i$  then we denote it  $\overrightarrow{E_{v_i}}$ ,  $v \in V$  since  $A_i = U_{i \in I}(\bigcap_{j=1}^n \overrightarrow{E_{v_i}})$  and  $B_i = U_{k \in I}(\bigcap_{j=1}^n \overrightarrow{E_{v_i}})$  then  $\overrightarrow{E_{v_i}}$  one of these Posts. Therefore  $A_i \cap B_i \in \tau_{ID}$ .

**Example 3.3:** Let  $D = (V, E)$  be digraph as in Figure (1), such that  $V = \{v_1, v_2, v_3, v_4, v_5, v_6\}$ ,  $E = \{e_1, e_2, e_3, e_4, e_5, e_6\}$ .

We have,  $\overrightarrow{E_{v_1}} = \emptyset$ ,  $\overrightarrow{E_{v_2}} = \{e_1\}$ ,  $\overrightarrow{E_{v_3}} = \{e_2\}$ ,  $\overrightarrow{E_{v_4}} = \{e_3\}$ ,  $\overrightarrow{E_{v_5}} = \{e_4, e_6\}$ ,  $\overrightarrow{E_{v_6}} = \{e_5\}$ . and  $S_{ID} = \{\emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4, e_6\}\}$ .

By taking finitely intersection the basis obtained is :  $\{\emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4, e_6\}\}$ . Then by taking all unions the topology can be written as:

$\tau_{ID} = \{ \emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4, e_6\}, \{e_1, e_2\}, \{e_1, e_3\}, \{e_1, e_5\}, \{e_2, e_3\}, \{e_2, e_5\}, \{e_3, e_5\}, \{e_1, e_2, e_3\}, \{e_1, e_2, e_5\}, \{e_1, e_3, e_5\}, \{e_2, e_3, e_5\}, \{e_1, e_4, e_6\}, \{e_2, e_4, e_6\}, \{e_3, e_4, e_6\}, \{e_5, e_4, e_6\}, \{e_1, e_2, e_3, e_5\}, \{e_1, e_2, e_4, e_6\}, \{e_1, e_3, e_4, e_5\}, \{e_1, e_5, e_4, e_5\}, \{e_2, e_3, e_4, e_5\}, \{e_2, e_5, e_4, e_5\}, \{e_3, e_5, e_4, e_5\}, \{e_1, e_2, e_3, e_4, e_6\}, \{e_1, e_2, e_5, e_4, e_6\}, \{e_1, e_3, e_5, e_4, e_6\}, \{e_2, e_3, e_5, e_4, e_6\} \}$ .

Then  $\tau_{ID}$  is topology is called digtopology  $\tau_{ID}$ .

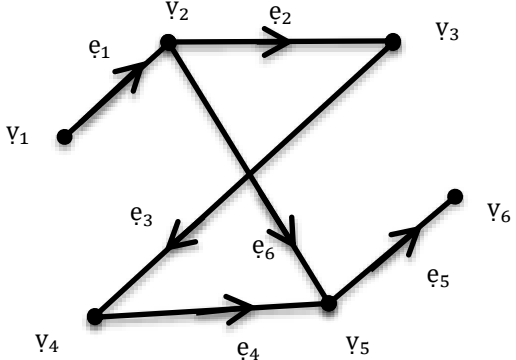


Fig.(1)

**Remark 3.4:** Let  $C_n$  be cyclic digraph if every edges are in the same directed then we get the digtopology  $\tau_{ID}$  on  $C_n$  is discrete, and if the edges are not all in the same direction we get the digtopology  $\tau_{ID}$  on  $C_n$  is not discrete.

This Remark illustrates in the next two Examples.

**Example 3.5:** Let  $C_5$  be cyclic digraph such that every edges are in the same direction, show in Figure (2).

We have,  $\overrightarrow{E_{v_1}} = \{e_2\}$ ,  $\overrightarrow{E_{v_2}} = \{e_3\}$ ,  $\overrightarrow{E_{v_3}} = \{e_4\}$ ,  $\overrightarrow{E_{v_4}} = \{e_5\}$ ,  $\overrightarrow{E_{v_5}} = \{e_6\}$ ,  $\overrightarrow{E_{v_6}} = \{e_1\}$ .

And  $S_{ID} = \{\{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}, \{e_5\}, \{e_6\}\}$ .  
 $\tau_{ID} = \{ \emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}, \{e_5\}, \{e_6\}, \{e_1, e_2\}, \{e_1, e_3\}, \{e_1, e_4\}, \{e_1, e_5\}, \{e_1, e_6\}, \{e_2, e_3\}, \{e_2, e_4\}, \{e_2, e_5\}, \{e_2, e_6\}, \{e_3, e_4\}, \{e_3, e_5\}, \{e_3, e_6\}, \{e_4, e_5\}, \{e_4, e_6\}, \{e_5, e_6\}, \{e_1, e_2, e_3\}, \{e_1, e_2, e_4\}, \{e_1, e_2, e_5\}, \{e_1, e_2, e_6\}, \{e_1, e_3, e_4\}, \{e_1, e_3, e_5\}, \{e_1, e_3, e_6\}, \{e_1, e_4, e_5\}, \{e_1, e_4, e_6\}, \{e_2, e_3, e_4\}, \{e_2, e_3, e_5\}, \{e_2, e_3, e_6\}, \{e_2, e_4, e_5\}, \{e_2, e_4, e_6\}, \{e_3, e_4, e_5\}, \{e_3, e_4, e_6\}, \{e_1, e_2, e_3, e_4\}, \{e_1, e_2, e_3, e_5\}, \{e_1, e_2, e_3, e_6\}, \{e_1, e_2, e_4, e_5\}, \{e_1, e_2, e_4, e_6\}, \{e_1, e_3, e_4, e_5\}, \{e_1, e_3, e_4, e_6\}, \{e_1, e_2, e_3, e_4, e_5\}, \{e_1, e_2, e_3, e_4, e_6\}, \{e_1, e_2, e_3, e_5, e_6\}, \{e_1, e_3, e_4, e_5, e_6\}, \{e_2, e_3, e_4, e_5, e_6\} \}$ . Then we get the digtopology  $\tau_{ID}$  of  $C_6$  is discrete topology.

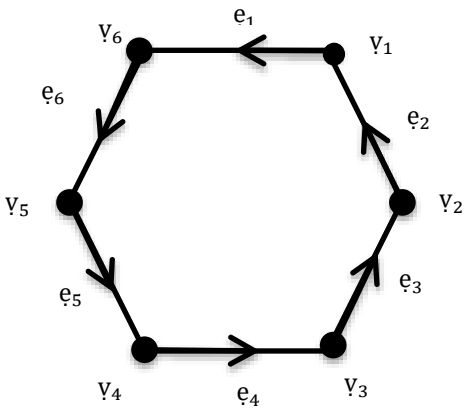
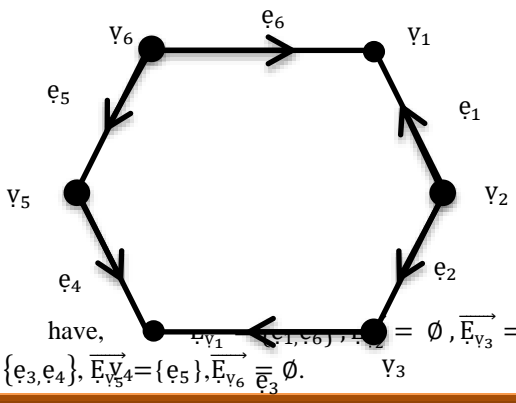


Fig.

**Example 3.6:** Let  $C_5$  be cyclic digraph such that the edges are not all in the same direction, show in Figure (3).



We have,  $\overrightarrow{E_{v_1}} = \{e_1, e_6\}$ ,  $\overrightarrow{E_{v_2}} = \{e_2\}$ ,  $\overrightarrow{E_{v_3}} = \{e_3\}$ ,  $\overrightarrow{E_{v_4}} = \{e_4\}$ ,  $\overrightarrow{E_{v_5}} = \{e_5\}$ ,  $\overrightarrow{E_{v_6}} = \{e_6\}$ .

Fig. (3)

And  $S_{TD} = \{\emptyset, \{e_1, e_6\}, \{e_3, e_4\}, \{e_2\}, \{e_5\}\}$ .

$\tau_{TD} = \{E(D), \emptyset, \{e_1, e_6\}, \{e_3, e_4\}, \{e_2\}, \{e_5\}, \{e_2, e_5\}, \{e_1, e_6, e_2\}, \{e_1, e_6, e_5\}, \{e_3, e_4, e_2\}, \{e_3, e_4, e_5\}, \{e_1, e_6, e_3, e_4\}, \{e_1, e_6, e_2, e_5\}, \{e_1, e_6, e_5, e_2\}, \{e_3, e_4, e_2, e_5\}, \{e_3, e_4, e_5, e_2\}, \{e_1, e_6, e_3, e_4, e_2\}, \{e_1, e_6, e_3, e_4, e_5\}\}$ . Then we get the digtopology  $\tau_{TD}$  of  $C_6$  is not discrete topology.

**Remark 3.7:** Let  $P_n$  be a path digraph if every edges are in the same directed then we get the digtopology  $\tau_{TD}$  on  $P_n$  is discrete, and if the edges are not all in the same direction we get the digtopology  $\tau_{TD}$  on  $P_n$  is not necessary discrete. This Remark illustrates in the next Examples.

**Example 3.8:** Let  $P_5$  be a path digraph, show in Figure (4).

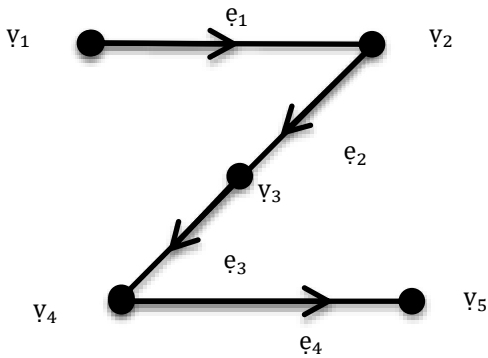


Fig.(4)

We have  $\overrightarrow{E_{v_1}} = \emptyset, \overrightarrow{E_{v_2}} = \{e_1\}, \overrightarrow{E_{v_3}} = \{e_2\}, \overrightarrow{E_{v_4}} = \{e_3\}, \overrightarrow{E_{v_5}} = \{e_4\}$ . And  $S_{TD} = \{\emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}\}$ .

$\tau_{TD} = \{E(D), \emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}, \{e_1, e_2\}, \{e_1, e_3\}, \{e_1, e_4\}, \{e_2, e_3\}, \{e_2, e_4\}, \{e_3, e_4\}, \{e_1, e_2, v_3\}, \{e_1, e_2, e_4\}, \{e_2, e_3, e_4\}\}$ .

We note the ine.to.digsp.  $\tau_{TD}$  is discret topology.

**Example 3.9:** Let  $P_5$  be a path digraph, show in Figure (5).

We have,  $\overrightarrow{E_{v_1}} = \{e_1\}, \overrightarrow{E_{v_2}} = \emptyset, \overrightarrow{E_{v_3}} = \{e_2\}, \overrightarrow{E_{v_4}} = \{e_3\}, \overrightarrow{E_{v_5}} = \{e_4\}$ .

And  $S_{TD} = \{\emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}\}$ .  $\tau_{TD} = \{E(D), \emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}, \{e_1, e_2\}, \{e_1, e_3\}, \{e_1, e_4\}, \{e_2, e_3\}, \{e_2, e_4\}, \{e_3, e_4\}, \{e_1, e_2, v_3\}, \{e_1, e_2, e_4\}, \{e_2, e_3, e_4\}\}$ .

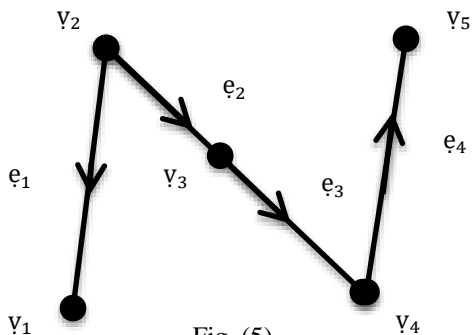


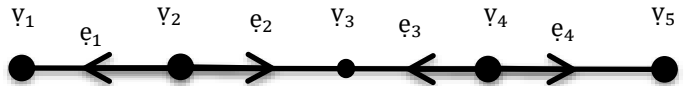
Fig. (5).

$\{e_1, e_2, e_3, e_4\}$ . We note the digtopology  $\tau_{TD}$  is discret topology.

**Example 3.10:** Let  $P_5$  be a path digraph, show in Figure (6) then we get the digtopology  $\tau_{TD}$  of  $P_5$  is not discrete.

We have,  $\overrightarrow{E_{v_1}} = \{e_1\}, \overrightarrow{E_{v_2}} = \emptyset, \overrightarrow{E_{v_3}} = \{e_2, e_3\}, \overrightarrow{E_{v_4}} = \emptyset, \overrightarrow{E_{v_5}} = \{e_4\}$ . And  $S_{TD} = \{\emptyset, \{e_1\}, \{e_2, e_3\}, \{e_4\}\}$ .

$\tau_{TD} = \{E(D), \emptyset, \{e_1\}, \{e_2, e_3\}, \{e_4\}, \{e_1, e_4\}, \{e_1, e_2, e_3\}, \{e_2, e_3, e_4\}, \{e_1, e_2, e_3, e_4\}\}$ . We note the ine.to.digsp.  $\tau_{TD}$  is not discret topology.



**Remark 3.11:** Let  $S_n$  be a star digraph then :

- (i) If every edges are direction to center vertex then the digtopology  $\tau_{TD}$  is indiscret topology.
- (ii) If every edges are not direction to center vertex then the digtopology  $\tau_{TD}$  is discret topology.
- (iii) If most edge is direction to center vertex then the digtopology  $\tau_{TD}$  is not discrete topology. This Remark illustrates in the next Examples.

**Example 3.12:** Let  $S_5$  be a star digraph, show in Figure(7).

In Figure(7)(A), every edges are directed to center vertex.

We have  $\overrightarrow{E_{v_1}} = \{e_1, e_2, e_3, e_4, e_5\}, \overrightarrow{E_{v_2}} = \emptyset, \overrightarrow{E_{v_3}} = \emptyset, \overrightarrow{E_{v_4}} = \emptyset, \overrightarrow{E_{v_5}} = \emptyset, \overrightarrow{E_{v_6}} = \emptyset$ . And  $S_{TD} = \{\emptyset, E(D)\}$

$\tau_{TD} = \{\emptyset, E(D)\}$  then we get the digtopology  $\tau_{TD}$  is indiscrete topology. But in Figure (7)(B) every the edges not direction to center vertex We note,  $\overrightarrow{E_{v_1}} = \emptyset, \overrightarrow{E_{v_2}} = \{e_1\}, \overrightarrow{E_{v_3}} = \{e_2\}, \overrightarrow{E_{v_4}} = \{e_3\}, \overrightarrow{E_{v_5}} = \{e_4\}, \overrightarrow{E_{v_6}} = \{e_5\}$ .

And  $S_{TD} = \{\emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}, \{e_5\}\}$

$\tau_{TD} = \{E(D), \emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}, \{e_5\}, \{e_1, e_2\}, \{e_1, e_3\}, \{e_1, e_4\}, \{e_1, e_5\}, \{e_2, e_3\}, \{e_2, e_4\}, \{e_2, e_5\}, \{e_3, e_4\}, \{e_3, e_5\}, \{e_4, e_5\}, \{e_1, e_2, e_3\}, \{e_1, e_2, e_4\}, \{e_1, e_2, e_5\}, \{e_1, e_3, e_4\}, \{e_1, e_3, e_5\}, \{e_1, e_4, e_5\}, \{e_2, e_3, e_4\}, \{e_2, e_3, e_5\}, \{e_2, e_4, e_5\}, \{e_3, e_4, e_5\}, \{e_1, e_2, e_3, e_4\}, \{e_1, e_2, e_3, e_5\}, \{e_1, e_2, e_4, e_5\}\}$ .

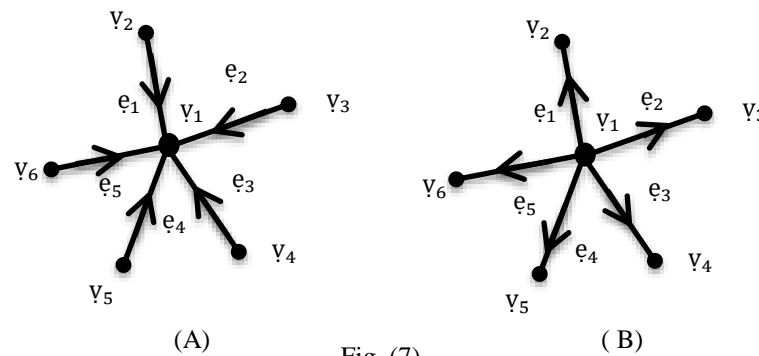


Fig. (7)

$\{e_1, e_3, e_4, e_5\}, \{e_2, e_3, e_4, e_5\}$ . we get the digtopology  $\tau_{TD}$  of  $S_5$  is discrete topology.

**Example 3.13:** Let  $S_4$  be a star digraph, show in Figure (8).

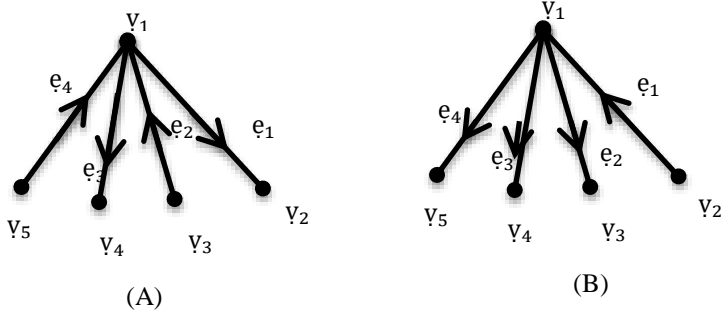


Fig. (8)

We note, in Figure (8)(A), is most edges is not same directed,  $\overrightarrow{E_{v_1}} = \{e_2, e_3, e_4\}$ ,  $\overrightarrow{E_{v_2}} = \{e_1\}$ ,  $\overrightarrow{E_{v_3}} = \emptyset$ ,  $\overrightarrow{E_{v_4}} = \{e_3\}$ ,  $\overrightarrow{E_{v_5}} = \emptyset$ . And  $S_{TD} = \{\emptyset, \{e_1\}, \{e_3\}, \{e_2, e_4\}\}$ .

$\tau_{TD} = \{E(D), \emptyset, \{e_1\}, \{e_3\}, \{e_2, e_3\}, \{e_1, e_3\}, \{e_1, e_2, e_4\}\}, \{e_3, e_2, e_4\}$   
 We get the digtopology  $\tau_{TD}$  is not discret topology .

But in Figure (8)(B), we note only one edges is directed to center vertexe , hence  $\overrightarrow{E_{v_1}} = \{e_1\}$ ,  $\overrightarrow{E_{v_2}} = \emptyset$ ,  $\overrightarrow{E_{v_3}} = \{e_2\}$ ,  $\overrightarrow{E_{v_4}} = \{e_3\}$ ,  $\overrightarrow{E_{v_5}} = \{e_4\}$ . And  $S_{TD} = \{\emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}\}$ .

$\tau_{TD} = \{E(D), \emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}, \{e_1, e_2\}, \{e_1, e_3\}, \{e_1, e_4\}, \{e_2, e_3\}, \{e_2, e_4\}, \{e_3, e_4\}, \{e_1, e_2, e_3\}, \{e_1, e_2, e_4\}, \{e_1, e_3, e_4\}, \{e_2, e_3, e_4\}\}$ . We get the digtopology  $\tau_{TD}$  of  $S_5$  is discrete topology.

**Proposition 3.14:** Suppose that  $\tau_{TD}$  is the digtopology of the adigraph  $D = (V, E)$ , then  $\{e\} \in \tau_{TD}$  if  $\overrightarrow{I_e^v} \neq \overrightarrow{I_e^w}$  for all  $e \in E$

**Prove:** Let  $e \in E$  then  $\overrightarrow{I_e^v} = \{v\}$  for some  $v \in V$  and by hypothesis  $\overrightarrow{I_e^v} \neq \overrightarrow{I_e^w}$  for all  $e \in E$  then we get  $e$  is only edge is directed to  $v$  and hence then  $\overrightarrow{E_v} = \{e\}$  and by definition digtopology  $\tau_{TD}$  we get  $\{e\} \in \tau_{TD}$ .

**Remark 3.15:** Let  $D = (V, E)$  be a digraph, then the digtopology  $\tau_{TD}$  is not necessary to be discrete topology in general .The Example 3.3, illustrate Remark 3.15 .

**Corollary 3.16:** Let  $D = (V, E)$  be a digraph if  $\overrightarrow{I_e^v} \neq \overrightarrow{I_e^w}$  for every distance pair of edge  $e, e \in E$  , then digtopology  $\tau_{TD}$  is discrete topology.

**Prove:** Since  $\overrightarrow{I_e^v} \neq \overrightarrow{I_e^w}$  for every distance pair of edge in digraph  $D$  then by Proposition 3.14,  $\{e\} \in \tau_{TD}$  for all  $e \in E$ , hence we get the digtopology  $\tau_{TD}$  is discrete topology.

**Example 3.17:** According to Example 3.5, we note that  $\overrightarrow{I_{e_1}^v} = \{v_5\}$ ,  $\overrightarrow{I_{e_2}^v} = \{v_1\}$ ,  $\overrightarrow{I_{e_3}^v} = \{v_2\}$ ,  $\overrightarrow{I_{e_4}^v} = \{v_3\}$ ,  $\overrightarrow{I_{e_5}^v} = \{v_4\}$  i.e  $\overrightarrow{I_e^v} \neq \overrightarrow{I_e^w} \forall e, e \in E$ , hence digtopology  $\tau_{TD}$  is discrete topology.

**Remark 3.18:** If  $D = (V, E)$  be reflexive digraph then digtopology  $\tau_{TD}$  is not discrete topology .

The Example illustrate Remark 3.18 .

**Example 3.19:** Let  $D = (V, E)$  be reflexive digraph show in Figure(9).

We note,  $\overrightarrow{E_{v_1}} = \{e_1\}$ ,  $\overrightarrow{E_{v_2}} = \{e_2, e_3\}$ ,  $\overrightarrow{E_{v_3}} = \{e_4, e_5\}$  .  
 And  $S_{TD} = \{\{e_1\}, \{e_2, e_3\}, \{e_4, e_5\}\}$ .  $\tau_{TD} = \{E(D), \emptyset, \{e_1\}, \{e_2, e_3\}, \{e_4, e_5\}, \{e_1, e_2, e_3\}\}$ ,

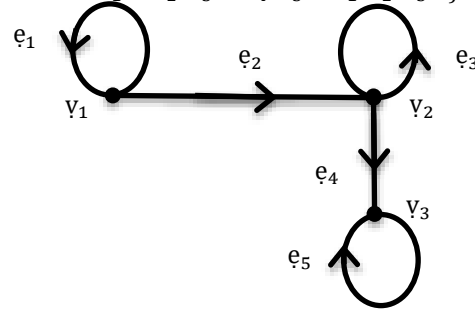


Fig.(9).

$\{e_1, e_4, e_5\}, \{e_2, e_3, e_4, e_5\}$ . We get the digtopology  $\tau_{TD}$  is not discrete topology.

**Proposition 3.20:** Let  $D = (V, E)$  be reflexive digraph and  $d(v) \leq 2$  then the digtopology  $\tau_{TD}$  is discrete topology.

**Prove :** Since the  $D = (V, E)$  reflexive digraph and  $d(v) \leq 2$  for all  $v \in V$  there exist only loop on every vertex and hence  $\forall v \in V$  we get  $E_v = \{e\}$ , where  $e = (v, v)$  and by definition digtopology  $\tau_{TD}$  implies  $\{e\} \in \tau_{TD}$  for all  $e \in E$ , thus digtopology  $\tau_{TD}$  is discrete topology. The Example illustrate proposition 3.20.

**Example 3.21:** Let  $D = (V, E)$  be reflexive digraph, show in Figure (10).

We note,  $\overrightarrow{E_{v_1}} = \{e_1\}$ ,  $\overrightarrow{E_{v_2}} = \{e_2\}$ ,  $\overrightarrow{E_{v_3}} = \{e_3\}$ ,  $\overrightarrow{E_{v_4}} = \{e_4\}$  .  
 And  $S_{TD} = \{\{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}\}$ .

$\tau_{TD} = \{E(D), \emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}, \{e_1, e_2\}, \{e_1, e_3\}, \{e_1, e_4\}, \{e_2, e_3\}, \{e_2, e_4\}, \{e_3, e_4\}, \{e_1, e_2, e_3\}, \{e_1, e_2, e_4\}, \{e_1, e_3, e_4\}, \{e_2, e_3, e_4\}\}$   
 We get the digtopology  $\tau_{TD}$  is discrete topology.

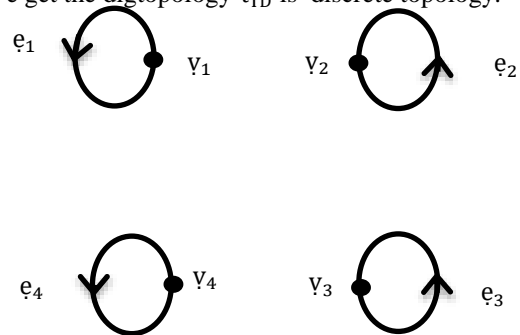


Fig.(10).

**Remark 3.22:** Let  $D = (V, E)$  be a symmetric digraph, then digtopology  $\tau_{TD}$  not necessary to be discrete topology in general .the following Example shows Remark 3.22.

**Example 3.23:** Let  $D = (V, E)$  be symmetric digraph in Figure (11) such that  $V = \{v_1, v_2, v_3\}$  ,  $E = \{e_1, e_2, e_3, e_4\}$

We have,  $\overrightarrow{E}_{v_1} = \{e_1\}$ ,  $\overrightarrow{E}_{v_2} = \{e_2, e_3\}$ ,  $\overrightarrow{E}_{v_3} = \{e_4\}$ .  
 $S_{TD} = \{\{e_1\}, \{e_4\}, \{e_2, e_3\}\}$ .  
 $\tau_{TD} = \{ \overrightarrow{E}(\mathcal{D}), \emptyset, \{e_1\}, \{e_4\}, \{e_2, e_3\}, \{e_1, e_4\}, \{e_1, e_2, e_3\}, \{e_4, e_2, e_3\} \}$ . We note the digtopology  $\tau_{TD}$  is not discrete topology.

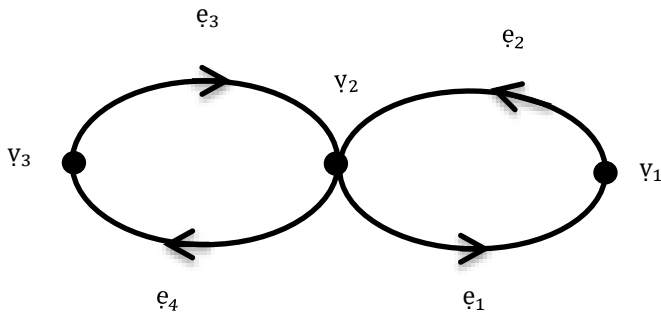


Fig.(11)

**Proposition 3.24:** If  $\mathcal{D} = (V, E)$  be a symmetric digraph and  $d(v) \leq 2$  for all  $v \in V$  then the digtopology  $\tau_{TD}$  is discrete topology.

**Prove:** Since  $\mathcal{D} = (V, E)$  be a symmetric digraph and  $d(v) \leq 2$  implies we get for all  $v \in V$  there exist at most one edge is directed to  $v$  for all  $v \in V$  [because if there exist two edge are directed to  $v$  and  $d(v) \leq 2$  implies  $\mathcal{D} = (V, E)$  is not symmetric] and hence,  $\overrightarrow{I}_e^v \neq \overrightarrow{I}_e^v$  for all  $e \in E$ , by proposition 3.14 then the digtopology  $\tau_{TD}$  is discrete topology.

the following Example shows the Proposition 3.24.

**Example 3.25:** Let  $\mathcal{D} = (V, E)$  be symmetric digraph in Figure (12) such that  $V = \{v_1, v_2, v_3, v_4, v_5\}$ ,  $E = \{e_1, e_2, e_3, e_4, e_5\}$ .

We note,  $\overrightarrow{E}_{v_1} = \{e_1\}$ ,  $\overrightarrow{E}_{v_2} = \{e_2\}$ ,  $\overrightarrow{E}_{v_3} = \{e_3\}$ ,  $\overrightarrow{E}_{v_4} = \{e_4\}$ ,  $\overrightarrow{E}_{v_5} = \{e_5\}$ . And  $S_{TD} = \{\{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}, \{e_5\}\}$ .  
 $\tau_{TD} = \{ \overrightarrow{E}(\mathcal{D}), \emptyset, \{e_1\}, \{e_2\}, \{e_3\}, \{e_4\}, \{e_5\}, \{e_1, e_2\}, \{e_1, e_3\}, \{e_1, e_4\}, \{e_1, e_5\}, \{e_2, e_3\}, \{e_2, e_4\}, \{e_2, e_5\}, \{e_3, e_4\}, \{e_3, e_5\}, \{e_4, e_5\}, \{e_1, e_2, e_3\}, \{e_1, e_2, e_4\}, \{e_1, e_2, e_5\}, \{e_1, e_3, e_4\}, \{e_1, e_3, e_5\}, \{e_1, e_4, e_5\}, \{e_2, e_3, e_4\}, \{e_2, e_3, e_5\}, \{e_2, e_4, e_5\}, \{e_3, e_4, e_5\}, \{e_1, e_2, e_3, e_4\}, \{e_1, e_2, e_3, e_5\}, \{e_1, e_2, e_4, e_5\}, \{e_1, e_3, e_4, e_5\}, \{e_2, e_3, e_4, e_5\} \}$ . We get the digtopology  $\tau_{TD}$  is discrete topology.

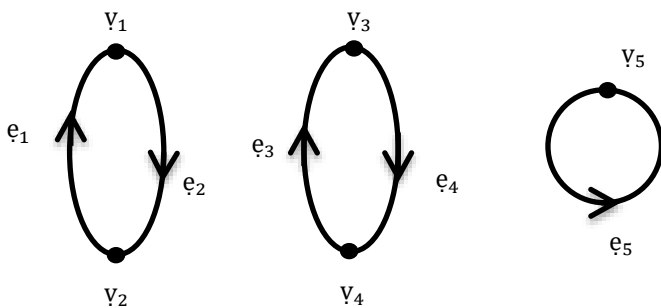


Fig. (12)

**Proposition 3.26:** The digtopology  $(E, \tau_{TD})$  of digraph  $\mathcal{D} = (V, E)$  is Alexandroff space.

**Prove:** It is adequate to show that arbitrary intersection of elements of  $S_{TD}$  is open, Let  $A \subseteq V$  :  $\bigcap_{v \in A} \overrightarrow{E}_v =$

$$\begin{cases} \overrightarrow{E}_v & \text{if } A \text{ contain one vertex } v \\ \emptyset & \text{otherwise} \end{cases}$$

And by Definition 3.1 of digtopology  $\tau_{TD}$  we get  $\emptyset, \overrightarrow{E}_v \in \tau_{TD}$ , then  $\bigcap_{v \in A} \overrightarrow{E}_v$  is open. Hence the digtopology  $\tau_{TD}$  is satisfies property of Alexandroff.

**Definition 3.27:** In any digraph  $\mathcal{D} = (V, E)$  since  $(E, \tau_{TD})$  is Alexandroff space, for each  $e \in E$ , the intersection of all open set containing  $e$  is the smallest open set containing  $e$  and denoted by  $U_e$ . Also the family  $M_{\mathcal{D}} = \{U_e | e \in E\}$  is the minimal basis for the digtopology  $\tau_{TD}$ .

**Proposition 3.28:** In any digraph  $\mathcal{D} = (V, E)$ ,  $U_e = \overrightarrow{E}_v$  where  $\overrightarrow{I}_e^v = \{v\}$  for every  $e \in E$ .

**Prove:**  $\Rightarrow$  Since every  $e \in E$  then  $\overrightarrow{I}_e^v = \{v\}$  for some  $v \in V$  and by Definition 3.1.1, of digtopology  $\tau_{TD}$ ,  $\overrightarrow{E}_v$  is open contain  $e$  and by Definition 3.27, of  $U_e$  then we get  $U_e \subseteq \overrightarrow{E}_v$ .

$\Leftarrow$  since  $U_e$  is open set and contain  $e$  then by definition digtopology  $\tau_{TD}$  there exist  $\overrightarrow{E}_v$  for some  $v \in V$  such that  $e \in \overrightarrow{E}_v \subseteq U_e$ , and since  $e \notin U_e$  implies  $e \notin \overrightarrow{E}_v$  when  $e \in \overrightarrow{E}_v$ , then  $e \notin \overrightarrow{E}_v$  when  $\overrightarrow{I}_e^v = \{v\}$  thus  $\overrightarrow{E}_v \subseteq U_e$  when  $\overrightarrow{I}_e^v = \{v\}$ , then  $U_e = \overrightarrow{E}_v$  where  $\overrightarrow{I}_e^v = \{v\}$  for every  $e \in E$ .

**Remark 3.29:** Let  $\mathcal{D} = (V, E)$  be a digraph, if  $\overrightarrow{I}_e^v \neq \overrightarrow{I}_e^v$  for all  $e \in E$  then  $U_e = \{e\}$ .

**Prove:** clear

**Theorem 3.30:** For any  $e, e' \in E$  in a digraph  $\mathcal{D} = (V, E)$  we have  $\overrightarrow{I}_e^v = \overrightarrow{I}_{e'}^v$  iff  $e \in U_{e'}$  i.e  $U_e = \{e' \in E | \overrightarrow{I}_e^v = \overrightarrow{I}_{e'}^v\}$ .

**Prove:**  $\Rightarrow$  Let  $\overrightarrow{I}_e^v = \overrightarrow{I}_{e'}^v$  to prove  $e \in U_{e'}$ . Since  $\overrightarrow{I}_e^v = \overrightarrow{I}_{e'}^v = \{v\}$ , implies  $e, e' \in \overrightarrow{E}_v$  and by Properties 3.28  $U_e = \overrightarrow{E}_v$  and  $U_{e'} = \overrightarrow{E}_v$ , and hence  $U_e = U_{e'} = \overrightarrow{E}_v$  then  $e \in U_{e'}$ .

$\Leftarrow$  let  $e \in U_{e'}$  and by Proposition 3.28, we get  $U_e = \overrightarrow{E}_v$  where  $\overrightarrow{I}_e^v = \{v\}$ , since  $e \in U_{e'}$  we get  $e \in \overrightarrow{E}_v$  where  $\overrightarrow{I}_{e'}^v = \overrightarrow{I}_e^v = \{v\}$  and hence  $\overrightarrow{I}_e^v = \overrightarrow{I}_{e'}^v$ .

**Remark 3.31:** The digtopology  $\tau_{TD}$  in any digraph  $\mathcal{D} = (V, E)$  is not necessary  $T_0$  in general.

**Example 3.32:** According to Example 3.12 we get the digtopology  $\tau_{TD} = \{\emptyset, \overrightarrow{E}(\mathcal{D})\}$  is not  $T_0$  because  $e_2, e_3 \in \overrightarrow{E}(\mathcal{D})$  But  $\nexists$  open set  $A$  such that  $e_2 \in A$  and  $e_3 \notin A$  or  $e_3 \in A$  and  $e_2 \notin A$ .

**Remark 3.33:** Let  $C_n$  be a cyclic digraph such that every edges are in the same direction then we get the digtopology  $\tau_{TD}$  is  $T_0$  and if the edges are not all in the same direction we get the the digtopology  $\tau_{TD}$  is not  $T_0$ .

The next Example are illustrates the remark 3.33.

**Example 3.34:**

- (i) By according Example 3.5 we note that the cyclic digraph  $C_5$  all edges in the same direction

and hence we note the digtopology  $\tau_{1D}$  on  $C_5$  is  $T_0$ .

- (ii) By according example 3.6 we note that the cyclic digraph  $C_6$  all edges are not in the same direction and hence we note the digtopology  $\tau_{1D}$  is not  $T_0$ , because  $e_3, e_4 \in E$  but  $\nexists A \in \tau_{1D}$  such that  $e_3 \in A$  and  $e_4 \notin A$  or  $e_3 \notin A$  and  $e_4 \in A$ .

**Remark 3.35 :** Let  $P_n$  be a path digraph such that every edges are in the same direction then we get the digtopology  $\tau_{1D}$  is  $T_0$ , and if the edges are not all in the same direction we get the digtopology  $\tau_{1D}$  is not necessary  $T_0$ .

The next two Example are illustrates the Remark 3.35 .

**Example 3.36:**

- (i) By according Example 3.8 we note that the path digraph  $P_5$  all edges in the same direction and hence we note the digtopology  $\tau_{1D}$  on  $P_5$  is  $T_0$ .
- (ii) By according Example 3.9 we note that the path digraph  $P_5$  all edges are not in the same direction and hence we note the digtopology  $\tau_{1D}$  is  $T_0$ .
- (iii) By according Example 3.10 we note that the path digraph  $P_5$  all edges are not in the same direction and hence we note the digtopology  $\tau_{1D}$  is not  $T_0$ , because  $e_2, e_3 \in E$  but  $\nexists A \in \tau_{1D}$  such that  $e_2 \in A$  and  $e_3 \notin A$  or  $e_3 \notin A$  and  $e_2 \in A$ .

**Remark 3.37:** Let  $S_n$  be a star digraph such that every edges are indgree to the center vertex then we get the digtopology  $\tau_{1D}$  is not  $T_0$ , and if the every edges are not indgree to the center vertex we get the digtopology  $\tau_{1D}$  is  $T_0$ . Also, if only one edge from set edges is directed in to center vertex then we have the digtopology  $\tau_{1D}$  is  $T_0$ .

The next Example are illustrates the Remark 3.37.

**Example 3.38 :**

- (i) By according Example 3.12 in Figure (9)(A) we note that the star digraph  $S_5$  every edges are indgree to the center vertex and hence we note the digtopology  $\tau_{1D}$  is not  $T_0$ , because  $e_2, e_4 \in E$  but  $\nexists A \in \tau_{1D}$  such that  $e_2 \in A$  and  $e_4 \notin A$  or  $e_2 \notin A$  and  $e_4 \in A$ .
- (ii) By according Example 3.12 in figure (7)(B) we note that the star digraph  $S_5$  every edges are not indgree to the center vertex we get the digtopology  $\tau_{1D}$  is  $T_0$ .

**Proposition 3.39:** The digtopology  $\tau_{1D}$  in any digraph  $D = (V, E)$  is  $T_0$  if and only if  $\vec{I}_e \neq \vec{I}_e$  for every distinct pair of edges  $e, e \in E$ .

**Prove:**  $\implies$  Suppos the digtopology  $\tau_{1D}$  is  $T_0$  to prove  $\vec{I}_e \neq \vec{I}_e$  for every distinct pair of edges  $e, e \in E$ .

If  $\vec{I}_e = \vec{I}_e$  then by Theorem 3.30  $e \in U_e$  and we get there exist  $u$  is open set such that  $e \in u$  and  $e \notin u$  implies the

digtopology  $\tau_{1D}$  is not  $T_0$  this contradiction .then  $\vec{I}_e \neq \vec{I}_e$  for every distinct pair of edges  $e, e \in E$ .

$\Leftarrow$  since  $\vec{I}_e \neq \vec{I}_e$  for every distinct pair of edges then by corollary 3.16 digtopology  $\tau_{1D}$  is discrete implies the digtopology  $\tau_{1D}$  is  $T_0$ .

The next Example is illusteted this Proposition 3.39.

**Example 3.40:** Acorroding Example 3.5, we note that  $\vec{I}_e \neq \vec{I}_e$  for every distinct pair of edges [since  $\vec{I}_{e_1} = \{v_5\}, \vec{I}_{e_2} = \{v_1\}, \vec{I}_{e_3} = \{v_2\}, \vec{I}_{e_4} = \{v_3\}, \vec{I}_{e_5} = \{v_4\}$ ] then we get the digtopology  $\tau_{1D}$  is  $T_0$ .

And we note, that in the Example 3.6  $\vec{I}_{e_1} = \{v_1\}, \vec{I}_{e_2} = \{v_3\}, \vec{I}_{e_3} = \{v_4\}, \vec{I}_{e_4} = \{v_4\}, \vec{I}_{e_5} = \{v_5\}, \vec{I}_{e_6} = \{v_1\}$  there exists  $\vec{I}_{e_1} = \vec{I}_{e_6}$  and hence, the digtopology  $\tau_{1D}$  is not  $T_0$ .

**Corollary 3.41:** The digtopology  $\tau_{1D}$  in any digraph  $D = (V, E)$  is  $T_0$  if and only if it is discrete .

**Prove :** The proof is easy by Properties 3.39 and Corollary 3.16.

**Remark 3.42:** The digtopology  $\tau_{1D}$  in any digraph  $D = (V, E)$  is not necessary  $T_1$  in general .

**Example 3.43:** According to Example 3.3 we get the digtopology  $\tau_{1D}$  is not  $T_1$  because  $e_4, e_6 \in E(D)$  But  $\nexists$  open set  $A$  such tha  $e_4 \in A$  and  $e_6 \notin A$  and  $e_6 \in A$  and  $e_4 \notin A$ .

**Remark 3.44:** Let  $C_n$  be a cyclic digraph such that every edges are in the same direction then we get the digtopology  $\tau_{1D}$  is  $T_1$ , and if the edges are not all in the same direction we get the digtopology  $\tau_{1D}$  is not  $T_1$ .

The next Example are illustrates the Remark 3.44.

**Example 3.45 :**

- (i) By according Example 3.5 we note that the cyclic digraph  $C_5$  all edges in the same direction and hence we note the digtopology  $\tau_{1D}$  on  $C_5$  is  $T_1$ .
- (ii) By according Example 3.6 we note that the cyclic digraph  $C_6$  all edges are not in the same direction and hence we note the digtopology  $\tau_{1D}$  is not  $T_1$ , because  $e_1, e_6 \in E$  but  $\nexists A \in \tau_{1D}$  such that  $e_1 \in A$  and  $e_6 \notin A$  and  $e_6 \in A$  and  $e_1 \notin A$ .

**Remark 3.46 :** Let  $P_n$  be a path digraph such that every edges are in the same direction then we get the digtopology  $\tau_{1D}$  is  $T_1$ , and if the edges are not all in the same direction we get the digtopology  $\tau_{1D}$  is not necessary  $T_0$ .

This Remark illustrates in the next Example.

**Example 3.47:**

- (i) By according Example 3.8 we note that the path digraph  $P_5$  all edges in the same direction and hence we note the digtopology  $\tau_{1D}$  on  $P_5$  is  $T_1$ .
- (ii) By according Example 3.9 we note that the path digraph  $P_5$  all edges are not in the same direction and hence we note the digtopology  $\tau_{1D}$  is  $T_1$ .

- (iii) By according Example 3.10 we note that the path digraph  $P_5$  all edges are not in the same direction and hence we digtopology  $\tau_{TD}$  is not  $T_1$ , because  $e_2, e_3 \in E$  but  $\nexists A \in \tau_{TD}$  such that  $e_2 \in A$  and  $e_3 \notin A$  and  $e_3 \in A$  and  $e_2 \in A$ .

**Proposition 3.48:** The digtopology  $\tau_{TD}$  in any digraph  $D = (V, E)$  is  $T_1$  if and only if  $\overrightarrow{I_e} \neq \overrightarrow{I_{\acute{e}}}$  for every distinct pairs of edges  $e, \acute{e} \in E$ .

**Prove:**  $\Rightarrow$  Suppos the digtopology  $\tau_{TD}$  is  $T_1$  to prove  $\overrightarrow{I_e} \neq \overrightarrow{I_{\acute{e}}}$  for every distinct pair of edges  $e, \acute{e} \in E$ .

If  $\overrightarrow{I_e} = \overrightarrow{I_{\acute{e}}}$  then by Theorem 3.30  $e \in U_{\acute{e}}$  and we get  $u$  is open set such that  $\acute{e} \in u$  and  $e \in u$  implies the digtopology  $\tau_{TD}$  is not  $T_1$  this contradiction, thus  $\overrightarrow{I_e} \neq \overrightarrow{I_{\acute{e}}}$  for every distinct pair of edges  $e, \acute{e} \in E$ .

$\Leftarrow$  Let  $\overrightarrow{I_e} \neq \overrightarrow{I_{\acute{e}}}$  for every distinct pair of edge by Corollary 3.16 implies the digtopology  $\tau_{TD}$  is discrete and hence  $\tau_{TD}$  is  $T_1$ .

**Corollary 3.49:** The digtopology  $\tau_{TD}$  in any digraph  $D = (V, E)$  is  $T_1$  if and only if it is discrete.

**Prove :** The proof is easy by Properties 3.48 and Corollary 3.16.

**Proposition 3.50:** The digtopology  $\tau_{TD}$  in any digraph  $D = (V, E)$  is  $T_0$  if and only if  $T_1$ .

**Prove:** By Proposition 3.46 and Proposition 3.41.

**Corollary 3.51:** Let  $D = (V, E)$  be a digraph, For every  $e \in E$  we have  $\overline{U_e} = \overline{E_v}$  where  $\overrightarrow{I_e} = \{v\}$ .

**Prove:** Let  $e \in E$  by Proposition 3.28 we get  $U_e = \overline{E_v}$  where  $\overrightarrow{I_e} = \{v\}$  Therefor  $\overline{U_e} = \overline{E_v}$  where  $\overrightarrow{I_e} = \{v\}$ .

**Corollary 3.52:** Given a digraph  $D = (V, E)$ . for every  $e \in E$   $\{e\} \subseteq \overline{U_e} = \overline{E_v}$  where  $\overrightarrow{I_e} = \{v\}$ .

**Prove:** Let  $u \in \{e\}$  this implies  $U \cap \{e\} \neq \emptyset$  for all open set  $U$  containing  $e$ . since  $\{e\} \subseteq U_e$  this implies  $U \cap U_e \neq \emptyset$  for all open set  $U$  containing  $e$ . hence  $e \in \overline{U_e}$  and so,  $\{e\} \subseteq \overline{U_e}$ . then by Corollary 3.42,  $\overline{\{e\}} \subseteq \overline{U_e} = \overline{E_v}$  where  $\overrightarrow{I_e} = \{v\}$ .

**Corollary 3.53:** For any  $e, \acute{e} \in E$  in a digraph  $D = (V, E)$  we have,  $\acute{e} \in \overline{\{e\}}$  if and only if  $\overrightarrow{I_e} = \overrightarrow{I_{\acute{e}}}$ .

**Prove:**  $\acute{e} \in \overline{\{e\}} \Leftrightarrow U \cap \{e\} \neq \emptyset$  for all open set  $U$  containing  $e \Leftrightarrow e \in U_{\acute{e}} \Leftrightarrow \overrightarrow{I_e} = \overrightarrow{I_{\acute{e}}}$ , by corollary 3.28.

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